Cybermedia Center Seminar: 18 Oct.2017

Introduction of direct solvers for large linear system with hands-on tutorial

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sparse direct solver

A : sparse matrix, nonsingular,

nonzero entries are in symmetric position : structural symmetric $A = \Pi_S^T \Pi_N^T W L D U W \Pi_N \Pi_S$,

 Π_S, Π_N : permutation, W: scaling matrix(e.g. diagonal) Three phases of direct solver:

- ▶ symbolic factorization to define Π_S to minimize fill-in : ordering
- ► numeric factorization with pivoting Π_N to reduce numerical round-off error
- forward-backward substitution

remarks

- unsymmetric nonzero pattern will be included in symmetric nonzero pattern by adding zero value at the symmetric position.
- sparse matrix is obtained from discretization of partial differential equation by FDM, FVM, or FEM.
- direct solver can factorize matrix with 1M DOFs by a multi-core processor.
- alternative solver is iterative method: Krylov subspace method (CG, GMRES), multigrid method, but direct solver is most robust

State of the art : software for sparse direct solver

	Software	parallel env.	elimination strategy	data manag.	pivoting	kernel detection
-	UMFPACK		multi-frontal	static	yes	no
	SuperLU_MT	shared	super-nodal	dynamic	yes	no
	Pardiso	shared/dist.	super-nodal	dynamic	yes + $\sqrt{\varepsilon}$ -p.	. no
	SuperLU_DIST	distributed	super-nodal	static	no, $\sqrt{\varepsilon}$ -p.	no
	MUMP S	dist./shared	multi-frontal	dynamic	yes	yes
	Dissection	shared	multi-frontal	static	yes	yes

T. A. Davis, I. S. Duff. A combined unifrontal/multifrontal method for unsymmetric sparse matrices,

ACM Trans. Math. Software, 25 (1999), 1–20.

J. W. Demmel, S. C. Eisenstat, J. R. Gilbert, X. S. Li, J. W. H. Liu.

A supernodal approach to sparse partial pivoting,

SIAM J. Matrix Anal. Appl., 20 (1999), 720-755.

O. Schenk, K. Gärtner. Solving unsymmetric sparse systems of liner equations with PARDISO,

Future Generation of Computer Systems, 20 (2004), 475–487.

X. S. Li, J. W. Demmel. SuperLU_DIST : A scalable distributed-memory sparse direct solver for unsymmetric linear systems,

ACM Trans. Math. Software, 29 (2003), 110–140.

P. R. Amestoy, I. S. Duff, J.-Y. L'Execellent. Mutlifrontal parallel distributed symmetric and unsymmetric solvers,

Comput. Methods Appl. Mech. and Engrg, 184 (2000) 501–520.

A. Suzuki, F.-X. Roux, A dissection solver with kernel detection for symmetric finite element matrices on shared memory computers,

Int. J. Numer. Meth. in Engng, 100 (2014) 136–164.

Hands-on tutorial to deal with direct solver packages

- Pardiso in Intel MKL, ver.2017.0.2
- MUMPS ver.5.1.1
- Dissection ver.1.0.3

Superscalar/Vector multicore parallel computers in CMC.

- ► VCC: Intel Xeon E5-2670V2@2.5GHz: 10 cores×2, 64GB mem.
- SX-ACE : NEC SX-ACE@1.0GHz: 4 cores×1, 64GB mem. 256GFlop/s

sparse matrces on structural mechanics/fluid dynamics from https://sparse.tamu.edu/ and original one

± 1 1 ±		0	
matrix	n	nnz	characteristic
poisson3Db	85,623	237,4949	unsymmetric
CoupCons3D	416,800	22,322,336	unsymmetirc
F1	343,791	13,590,452	symmetric
stokes.skewsym	565,003	16,715,009	unsymmetirc, singular

commercial CeCILL-C GPLv3/CeCILL-C

100GFlop/s

development of multi-core CPU

- clock speed of CPU never becomes faster in recent ten years stagnating at around 3GHz
 - e.g., Intel CPUs

clock	#cores	name	energy	year
3.6 GHz	2	Pentium D 960	130W	2006
3.2 GHz	4	Xeon E7-8893V4	140W	2016
2.1 GHz	18	Xeon E5-2695V4	120W	2016
1.5 GHz	72	Xeon Phi 7290F	260W	2016

one core in a note $pc \simeq$ one core in super computers

- multi-core processor with more than 64 cores
 - *LDU*-factorization with multi-frontal ordering
 - \Rightarrow no change of discretization of PDE up to 1M DOF
- advanced code for floating point operations

BLAS **level 3**, DGEMM

ordering of sparse matrix

sparse matrix needs to be re-ordered

- to reduce fill-in
- to increase parallelization of factorization
- to increase size of block structure

 \rightarrow multi-front \rightarrow supernode



nested-dissection parallel computation

reverse Cuthill-McKee sequential computation

original matrix band matrix

example:

7 stencil of Poisson equation, 11^3 nodes.



nested dissection by graph decomposition



A. George. Numerical experiments using dissection methods to solve n by n grid problems. SIAM J. Num. Anal. 14 (1977),161–179.

software package:

METIS : V. Kumar, G. Karypis, A fast and high quality multilevel scheme for partitioning irregular graphs. SIAM J. Sci. Comput. 20 (1998) 359–392.

SCOTCH : F. Pellegrini J. Roman J, P. Amestoy, Hybridizing nested dissection and halo approximate minimum degree for efficient sparse matrix ordering. Concurrency: Pract. Exper. 12 (2000) 69–84.

- each leaf can be computed in parallel \leftarrow multi-front
- load unbalance? because of different size of subdomains
- parallel computation of higher levels? # cores > # subdomains

recursive generation of Schur complement



12 13 11

sparse part : completely in parallel dense part : better use of BLAS 3; dgemm, dtrsm

pivoting strategy

full pivoting : $A = \Pi_L^T L U \Pi_R$ find $\max_{k < i,j \le n} |A(i,j)|$



partial pivoting : $A = \Pi L U$ find $\max_{k < i \le n} |A(i, k)|$



symmetric pivoting : $A = \Pi^T L D U \Pi$ find max_{k < i \le n} |A(k, k)|



 $\begin{array}{l} 2 \times 2 \text{ pivoting : } A = \prod^{T} L \tilde{D} U \Pi \\ \text{find } \max_{k < i, j \le n} \text{det} \begin{vmatrix} A(i, i) & A(i, j) \\ A(j, i) & A(j, j) \end{vmatrix}$



sym. pivoting is mathematically not always possible

LDU factorization with rank-1 update

$$\begin{bmatrix} a_{11} & \beta_1^T \\ \alpha_1 & A_{22} \end{bmatrix} = \begin{bmatrix} 1 & 0 \\ \alpha_1 a_{11}^{-1} & S_{22} \end{bmatrix} \begin{bmatrix} a_{11} & \beta_1^T \\ 0 & I_2 \end{bmatrix}$$
$$= \begin{bmatrix} 1 & 0 \\ \alpha_1 a_{11}^{-1} & S_{22} \end{bmatrix} \begin{bmatrix} a_{11} & 0 \\ 0 & I_2 \end{bmatrix} \begin{bmatrix} 1 & a_{11}^{-1} \beta_1^T \\ 0 & I_2 \end{bmatrix}$$

Schur complement $S_{22} = A_{22} - \alpha_1 a_{11}^{-1} \beta_1^T$ rank-1 update by dger LDU-factorization algorithm with symmetric pivoting do $k = 1, \cdots, N$ find max|A(l, l)| with $k < l \leq n$, exchange rows and columns : $A(k, *) \leftrightarrow A(l, *), A(*, k) \leftrightarrow A(*, l)$. dscal A(k, j) / = A(k, k) $k < j \le N$, dscal A(i,k) / = A(k,k) $k < i \le N$, dger $A(i, j) = A(i, k)A(k, k)^{-1}A(k, j)$ k < i, j < N. kl

LDU factorization of positive matrix : 1/2

$$\begin{split} A: \operatorname{coercive} \Leftrightarrow & (A\,x,x) > 0 \; \forall x \neq 0 \\ \Leftrightarrow & \frac{A + A^T}{2} : \operatorname{sym. positive definite} \; \Rightarrow \; \exists A^{-1} \\ x^T Ax = (Ax,x) = \frac{1}{2} \left\{ (A\,x,x) + (A\,x,x) \right\} = \frac{1}{2} \left\{ (A\,x,x) + (x,A^T\,x) \right\} \\ &= \left(\frac{A + A^T}{2} x, x \right) \\ A = \begin{bmatrix} A_{11} & \beta_1 \\ \alpha_1^T & \alpha_{22} \end{bmatrix} : \operatorname{coercive} \Rightarrow \exists \; A_{11}^{-1}, \; s_{22} \neq 0 \\ & \begin{bmatrix} A_{11} & \beta_1 \\ \alpha_1^T & \alpha_{22} \end{bmatrix} = \begin{bmatrix} I_1 & 0 \\ \alpha_1^T A_{11}^{-1} & 1 \end{bmatrix} \begin{bmatrix} A_{11} & \beta_1 \\ 0 & s_{22} \end{bmatrix} \\ & s_{22} = \alpha_{22} - \alpha_1^T A_{11}^{-1} \; \beta_1 \in \mathbb{R} \end{split}$$
proof
$$(A_{11}x_1, x_1) = \left(A \begin{bmatrix} x_1 \\ 0 \end{bmatrix}, \begin{bmatrix} x_1 \\ 0 \end{bmatrix} \right) > 0 \Rightarrow \exists A_{11}^{-1} \end{split}$$

$$\begin{bmatrix} A_{11} & \beta_1 \\ \alpha_1^T & \alpha_{22} \end{bmatrix} \begin{bmatrix} A_{11}^{-1} \beta_1 \\ -1 \end{bmatrix} = \begin{bmatrix} 0 \\ s_{22} \end{bmatrix}, \quad s_{22} = 0 \Rightarrow \mathsf{dimker} A \ge 1 \Rightarrow \Leftarrow$$

LDU factorization of positive matrix : 2/2

$$\begin{array}{ll} A: \text{coercive} \Leftrightarrow & (A\,x,x) > 0 \; \forall x \neq 0 \\ \\ \Leftrightarrow & \displaystyle \frac{A+A^T}{2}: \text{sym. positive definite} \end{array}$$

By recursively applying previous argument,

$$\begin{split} A &= \begin{bmatrix} A_{11} & \beta_1 \\ \alpha_1^T & \alpha_{22} \end{bmatrix} : \text{coercive} \Rightarrow \\ \begin{bmatrix} A_{11} & \beta_1 \\ \alpha_1^T & \alpha_{22} \end{bmatrix} = \begin{bmatrix} I_1 & 0 \\ \alpha_1^T A_{11}^{-1} & 1 \end{bmatrix} \begin{bmatrix} A_{11} & 0 \\ 0 & s_{22} \end{bmatrix} \begin{bmatrix} I_{11} & A_{11}^{-1} \beta_1 \\ 0 & 1 \end{bmatrix} \\ &= \begin{bmatrix} L_{11} & 0 \\ \alpha_1^T U_{11}^{-1} D_{11}^{-1} & 1 \end{bmatrix} \begin{bmatrix} D_{11} & 0 \\ 0 & s_{22} \end{bmatrix} \begin{bmatrix} U_{11} & D_{11}^{-1} L_{11}^{-1} \beta_1 \\ 0 & 1 \end{bmatrix} \end{split}$$

- ► By mathematically exact computation, coercive matrix *A* has *LDU*-factorization with any permutation.
- By floating point computation, numerical pivoting is necessary to get stable factorization with good accuracy.

2×2 pivot for indefinite matrix

symmetric matrix

$$\begin{bmatrix} \frac{1}{4} & \frac{5}{4} & \frac{1}{2} \\ \frac{5}{4} & \frac{1}{4} & \frac{1}{2} \\ \frac{1}{2} & \frac{1}{2} & 1 \end{bmatrix} = \begin{bmatrix} 1 & & \\ 5 & 1 & \\ 2 & \frac{1}{3} & 1 \end{bmatrix} \begin{bmatrix} \frac{1}{4} & & \\ & -6 & \\ & & \frac{2}{3} \end{bmatrix} \begin{bmatrix} 1 & 5 & 2 \\ 1 & \frac{1}{3} \\ & & 1 \end{bmatrix}$$

by symmetric permutation

$$\begin{bmatrix} 1 & \frac{1}{2} & \frac{1}{2} \\ \frac{1}{2} & \frac{1}{4} & \frac{5}{4} \\ \frac{1}{2} & \frac{5}{4} & \frac{1}{4} \end{bmatrix} = \begin{bmatrix} 1 & & \\ \frac{1}{2} & 1 & \\ \frac{1}{2} & 0 & 1 \end{bmatrix} \begin{bmatrix} 1 & & \\ 0 & 1 \\ & 1 & 0 \end{bmatrix} \begin{bmatrix} 1 & \frac{1}{2} & \frac{1}{2} \\ & 1 & 0 \\ & & 1 \end{bmatrix}$$

pivot strategy to take the largest entry may fail for indefinite matrix. a remedy is to use 2×2 pivot.

an algorithm to mix 1×1 and 2×2 pivots for sym. indefinite matrix:

J. R. Bunch, L. Kaufman. Some stable methods for calculating inertia and solving symmetric linear systems, *Math. Comput*, 31 (1977) 163–179.

$\sqrt{\varepsilon}$ -perturbation

a regularization technique

$$\begin{bmatrix} A_{11} & A_{12} \\ & 0 & \alpha \\ A_{21} & \beta & 0 \end{bmatrix} \rightarrow \begin{bmatrix} A_{11} & A_{12} \\ & \sqrt{\varepsilon} & \alpha \\ A_{21} & \beta & 0 \end{bmatrix}$$

- iterative refinement to improve accuracy of a solution
- user can/have to specify perturbation parameter for unsymmetric matrix (default = 10⁻¹³ for Pardiso)

symmetric pivoting with postponing for block strategy

 nested-dissection decomposition may produce singular sub-matrix for indefinite matrix

 τ : given threshold for postponing, 10^{-2} for <code>MUMPS</code>, <code>Dissection</code> $|A(i,i)|/|A(i-1,i-1)| < \tau \ \Rightarrow \{A(k,j)\}_{i \leq k,j}$ are postponed



Schur complement matrix from postponed pivots is computed kernel detection algorithm : QR-factorization by MUMPS

kernel detection (rank deficient problem)

$$\begin{bmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{bmatrix} = \begin{bmatrix} A_{11} & 0 \\ A_{21} & S_{22} \end{bmatrix} \begin{bmatrix} I_1 & A_{11}^{-1}A_{12} \\ 0 & I_2 \end{bmatrix} S_{22} = 0 \Rightarrow \operatorname{Ker} A = \begin{bmatrix} A_{11}^{-1}A_{12} \\ -I_2 \end{bmatrix}$$

symmetric semi-positive definite, m + k = 4 + 6 = 10by Householder-QR factorization:

4.60e-02	-1.20e-02	2.91e-03	1.16e-02	2.24e-02	-9.33e-05	-3.60e-02	8.22e-03	-7.77e-03	-2.90e-02
0.0	3.84e-02	4.84e-03	-2.21e-02	1.87e-02	1.30e-03	-9.14e-03	-2.74e-02	1.48e-02	-1.91e-02
0.0	0.0	2.96e-02	1.68e-03	-2.55e-02	1.11e-04	1.28e-02	-1.04e-04	-1.12e-03	-1.20e-02
0.0	0.0	0.0	1.28e-02	-1.66e-03	8.48e-04	-4.29e-05	-7.90e-04	-8.56e-03	2.51e-03
0.0	0.0	0.0	0.0	1.23e-11	-5.49e-13	-8.30e-12	1.67e-13	2.10e-14	-7.08e-12
0.0	0.0	0.0	0.0	0.0	6.70e-13	-1.02e-13	-5.33e-13	-1.62e-13	1.18e-13
0.0	0.0	0.0	0.0	0.0	0.0	3.33e-13	-2.48e-14	-6.61e-14	-3.18e-13
0.0	0.0	0.0	0.0	0.0	0.0	0.0	1.22e-13	4.46e-15	-4.34e-14
0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	3.05e-14	8.16e-15
0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	-1.09e-14

how to set threshold to distinguish between non-zero (1.28e-02) and zero (1.23e-11) values ?

Pardisono capability of kernel detection.MUMPSuser has to choose this value.Dissectionan algorithm by measuring dimension of residual of
matrix with a projection onto the image space.

kernel detection algorithm based on LDU (Dissection)

 $A: N \times N$ unsymmetric, dimKer $A = k \ge 1$, dimIm $A \ge m$. two parameters: l, n, which define size of factorization,

$$\begin{array}{c} N-n \uparrow \\ n \uparrow \end{array} \begin{bmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{bmatrix} = \begin{bmatrix} A_{11} & 0 \\ A_{21} & S_{22} \end{bmatrix} \begin{bmatrix} I_1 & A_{11}^{-1}A_{12} \\ 0 & I_2 \end{bmatrix} \quad \widetilde{\mathrm{Im}}_n = \mathrm{span} \begin{bmatrix} \widetilde{A_{11}^{-1}A_{12}} \\ -I_2 \end{bmatrix}^{\perp}$$

• projection :
$$P_n^{\perp} : \mathbb{R}^N \to \widetilde{\mathsf{Im}_n}$$

► solution in subspace,
$$\widetilde{A}_{N-l}^{\dagger} b = \begin{bmatrix} \widetilde{A_{11}^{-1}} b_1 \\ 0 \end{bmatrix}, b = \begin{bmatrix} b_1 \\ b_2 \end{bmatrix} \qquad \uparrow \ l \qquad l$$

 A_{11}^{-1} : computed in <u>quadruple-precision</u> with perturbation to simulate double-precision round-off error. *n*: candidate of dimension of the kernel compute for l = n - 1, n, n + 1

$$\begin{aligned} \operatorname{err}_{l}^{(n)} &:= \max \left\{ \max_{x=[0 \ x_{l}] \neq 0} \frac{||P_{n}^{\perp}(\tilde{A}_{N-l}^{\dagger}A \ x - x)||}{||x||}, \max_{x=[x_{N-l} \ 0] \neq 0} \frac{||\tilde{A}_{N-l}^{\dagger}A \ x - x||}{||x||} \right\} \\ n &= k + 1 \quad \Leftrightarrow \quad \operatorname{err}_{k}^{(k+1)} \approx 0 \quad \wedge \quad \operatorname{err}_{k+1}^{(k+1)} \approx 0 \quad \wedge \quad \operatorname{err}_{k+2}^{(k+1)} \sim 1 \\ n &= k \quad \Leftrightarrow \quad \operatorname{err}_{k-1}^{(k)} \gg 0 \quad \wedge \quad \operatorname{err}_{k}^{(k)} \approx 0 \quad \wedge \quad \operatorname{err}_{k+1}^{(k)} \sim 1 \\ n &= k - 1 \quad \Leftrightarrow \quad \operatorname{err}_{k-2}^{(k-1)} \gg 0 \quad \wedge \quad \operatorname{err}_{k-1}^{(k-1)} \gg 0 \quad \wedge \quad \operatorname{err}_{k}^{(k-1)} \sim 1 \end{aligned}$$

Theoretically $\neg A_{N-k+1}^{-1}$, but $\operatorname{err}_{k-1}^{(k)}$ is computable; $\sim ||A_N^{-1}A_N - I_N||$.

LDU factorization with rank-b update

$$\begin{bmatrix} A_{11} & A_{21} \\ A_{21} & A_{22} \end{bmatrix} = \begin{bmatrix} A_{11} & 0 \\ A_{21}A_{11}^{-1} & S_{22} \end{bmatrix} \begin{bmatrix} I_1 & A_{11}^{-1}A_{12} \\ 0 & I_2 \end{bmatrix}$$
$$= \begin{bmatrix} L_{11} & 0 \\ A_{21}U_{11}^{-1}D_{11}^{-1} & I_2 \end{bmatrix} \begin{bmatrix} D_{11} & 0 \\ 0 & S_{22} \end{bmatrix} \begin{bmatrix} U_{11} & D_{11}^{-1}L_{11}A_{12}^{-1} \\ 0 & I_2 \end{bmatrix}$$

Schur complement $S_{22} = A_{22} - A_{21}A_{11}^{-1}A_{12}$ rank-*b* update by dgemm *LDU*-factorization algorithm by block strategy do $k = 1, \dots, m = N/b$ *LDU*-factorization $A_{kk} = L_{kk}D_{kk}U_{kk}$ with symmetric pivoting. solve multiple-RHS $L_{kk}[\{X_j\}] = [A_{k+1}, \dots A_{km}],$ solve multiple-RHS $U_{kk}^T[\{Y_j^T\}] = [A_{k+1k}^T, \dots A_{mk}^T], j$ scaling $[\{Y_k\}] = D_{kk}[\{Y_k\}]$

k

i

LDU

 $A_{ik}U^{-T}$

 $L^{-1}A_{ki}$

rank-*b* update
$$[\{S_{ij}\}] = [\{A_{ij}\}] - [\{Y_j^T X_i\}].$$

- ► for modern CPU (multi-arithmetic units)
- parallelization
- ? limited search of pivots

task dependency analysis : example by dense matrix (Dissection)

LDL^T factorization $A_{11} A_{12} A_{13} A_{14} dtrsm + dcaling$ $\begin{vmatrix} A_{23} & A_{24} \\ A_{33} & A_{34} \end{vmatrix} \qquad \begin{vmatrix} A_{22} & A_{23} & A_{24} \\ A_{33} & A_{34} \end{vmatrix} - = \begin{vmatrix} A_{21} \\ A_{31} \end{vmatrix} \begin{vmatrix} A_{11}^{-1} \\ A_{11} \end{vmatrix} \begin{vmatrix} A_{12} & A_{13} & A_{14} \\ A_{31} \end{vmatrix}$ A₂₂ A₂₃ A₂₄ A₄₄ A₄₁ A44 daemm $\alpha^{(1)}$: factorization $L_{11}D_{11}L_{11}^T = A_{11}^{(1)}$ $\beta_{h}^{(1)}$: DTRSM + scaling $D_{11}^{-1}(L_{11}^{-1}A_{1h}^{(1)})$ $\gamma_{h,l}^{(1)}$: DGEMM $A_{k,l} = (L_{11}^{-1}A_{1,k})^T D_{11}^{-1} (L_{11}^{-1}A_{1,l}^{(1)})$ $\alpha^{(2)}$: factorization $L_{22}D_{22}L_{22}^T = A_{22}^{(2)}$ $\alpha^{(1)} \leftarrow \{\beta_2^{(1)}, \beta_2^{(1)}, \beta_4^{(1)}\} \leftarrow \{\gamma_{2,2}^{(1)}, \gamma_{2,2}^{(1)}, \gamma_{2,2}^{(1)}, \cdots, \gamma_{4,4}^{(1)}\} \leftarrow \alpha^{(2)} \leftarrow \{\beta_2^{(2)}, \beta_4^{(2)}\} \leftarrow \beta_{4,2}^{(2)}, \beta_{4,2}^{(2)}\} \leftarrow \beta_{4,2}^{(2)}, \beta_{4,2}^{(2)}, \beta_{4,2}^{(2)}, \beta_{4,2}^{(2)}\} \leftarrow \beta_{4,2}^{(2)}, \beta_{4,2}^{(2)}, \beta_{4,2}^{(2)}, \beta_{4,2}^{(2)}, \beta_{4,2}^{(2)}\} \leftarrow \beta_{4,2}^{(2)}, \beta$ $\{\gamma_{2,2}^{(2)}, \gamma_{2,4}^{(2)}, \gamma_{4,4}^{(2)}\} \leftarrow \alpha^{(3)} \cdots$ factorization is critical task, others depending on. 1 cf. DAG + runtime $\alpha^{(1)} \leftarrow \{\beta_2^{(1)}, \gamma_2^{(1)}, \alpha^{(2)}, \beta_3^{(1)}, \beta_4^{(1)}\} \leftarrow \{\gamma_2^{(1)}, \gamma_3^{(1)}, \cdots, \gamma_{4,4}^{(1)}\} \leftarrow$ $\{\beta_{2}^{(2)},\gamma_{2,2}^{(2)},\alpha_{3}^{(3)},\beta_{4}^{(2)}\} \leftarrow \{\gamma_{2,4}^{(2)},\gamma_{4,4}^{(2)}\} \leftarrow \beta_{2}^{(1)},\gamma_{2,2}^{(1)},\alpha_{3}^{(2)}\}$ atomic critical path is analyzed \Rightarrow asynchronous task execution

asynchronous task execution on multi-core (Dissection)



tasks of Schur complement in nested-dissection level 3, factorization in level 2 are scheduled together.

tasks on critical path are scheduled statically. others : 20% dynamic.

sparse matrix format : 1/3

```
n : # of rows
nnz : # of nonzeros
[A]<sub>ij</sub> : nonzero entries at (i, j)
COO (Coordinate) format MUMPS
int irow[nnz];
int jcol[nnz];
double coef[nnz];
```

 CSR (Compressed Sparse Row) / CRS (Compressed Row Storage) format Pradiso, Dissection

```
int ptrow[n+1];
int indcol[nnz];
double coef[nnz];
```

```
 [A]_{ij} = \operatorname{coef}[k] \\ j = \operatorname{indcol}[k], \operatorname{ptrow}[i] \le k < \operatorname{ptrow}[i+1]
```

sparse matrix format, zero-based index : 2/3

an example, 5×5 unsymmetric matrix, n = 5, nnz = 15.

									0	1	Ż	2	3 4	ł				
	1.1	1.2		1	.4			0	0	1		2	2					
	2.1	2.2	2.3	}		2.5		1	3	4	5	5	6	6				
		3.2	3.3	;				2		7	8	3						
	4.1			0	.0	4.5		3	9			1	0 1	1				
		5.2		5	.4	5.5		4		12	2	1	31	4				
	i	C)			1				2		3			4			5
ptr	ow[i] ()			3				7		9			12			15
ind	col[k] ()	1	3	0	1	2	4	1	2	0	3	4	1	3	4	
со	ef[k] 1	.1 1	.2	1.4	2.1	2.2	2.3	2.5	3.2	3.3	4.1	0.0	4.5	5.2	5.4	5.5	

diagonal entry should exist even if the value is 0

Pardiso, Dissection

indcol[] should be in ascending order in each row

sparse matrix format, zero-based index : 3/3

 5×5 symmetric matrix, upper triangular, n = 5, nnz = 10.

									0	1	- 2	3		4
	1.1	1.	2	1	1.4			0	0	1		2		
		2.	2 2	.3		2.5		1		3	4		!	5
			3	.3				2			6			
				(0.0	4.5		3				7		8
						5.5		4					ę	9
	i		0			1			2	3		4	5	
ptr	ow[i]	0			3			6	7		9	10	
ind	col[k]	0	1	3	1	2	4	2	3	4	4		
со	ef[k]	1.1	1.2	1.4	2.2	2.3	2.5	3.3	0.0	4.5	5.5		

- upper triangular matrix is accepted by Pardiso
- MUMPS accepts COO in upper or lower triangular sparse matrix with 1-based index in any order in the array of triplet

usage of Paridso from C/C++

```
MKL INT *ptrow = new MKL INT[n + 1]; // CSR data
MKL_INT *indcol = new MKL_INT[nnz];
double *coef = new double[nnz];
double *x = new double[n]: // solution
double *y = new double[n]; // RHS
void *pt[64]; // to keep internal pointers
MKL_INT *iparm = new MKL_INT[64]; // parameters!
MKL_INT mtype = 11; // structurally symmetric
MKL INT nrhs = 1;
MKL_INT phase;
MKL_INT maxfct = 1, mnum = 1, msglvl = 1, error;
MKL_INT idum; // dummy pointer instaed of user
                    // providing permutation
phase = 11; // symbolic factorization
pardiso(pt, &maxfct, &mnum, &mtype, &phase, &n,
        (void *)coef, ptrow, indcol, &idum, &nrhs,
        iparm, &msglvl, (void *)y, (void *)x,
        &error);
phase = 22; // numeric factorization
phase = 33; // Fw/Bw substitution
phase = -1; // free working data
```

usage of MUMPS from C/C++

```
MUMPS INT *irow = new MUMPS INT[nnz];
MUMPS_INT *jcol = new MUMPS_INT[nnz];
double *coef = new double[nnz];
double *x = new double[n]; // solution&RHS
DMUMPS_STRUC_C id;
id.job = (-1); // job init
id.par = 1;
id.sym = isSym ? 2 : 0;
id.comm fortran = USE COMM WORLD; // dummy MPI communicator
dmumps_c(&id);
id.job = 1; // symbolic facotrization
id.n = n; id.nz = nnz; id.irn = irow; id.jcn = jcol;
// id.icntl[] set parameters
dmumps c(&id);
id.job = 2; // numeric factorization
id.a = coef;
dmumps_c(&id);
id.job = 3; // Fw/Bw substitution
id.nrhs = 1;
id.rhs = x;
dmumps_c(&id);
id.job = -2; // free working data
```

usage of Dissection (fortran-interface) from C/C++

```
int *ptrow = new int[n + 1]; // CSR data
int *indcol = new int[nnz];
double *coef = new double[nnz];
double *x = new double[n]; // solution & RHS
uint64_t *dslv = new uint64_t; // Dissection solver object
int real_or_complex = 1; // 1 : real, 2 : complex
int verbose=1; // write messages
int call = 0;
diss_init(*dslv, call, real_or_complex, num_threads, verbose);
int sym = 0;
                 // unsymmetric, 1: symmetric upper
int ordering = 0; // 0 : SCOTCH, 1 : METIS
diss_s_fact(*dslv, n, ptrow, indcol, sym, ordering);
int scaling = 1; // diagonal scaling
int indefinte_flag = 1; // 0 for positive semi-definite
double eps_pivot = 1.0e-2; // threshold for postponing
diss_n_fact(*dslv, coef, scaling, eps_pivot, indefinite_lfag);
int projection = 1; // find solution in the image space
int transpose = 0; // 1 for A<sup>T</sup> x = b
diss_solv_1(*dslv, x, projectoin, transpose);
diss_free(*dslv); // clean the Dissection solver obj.
```

GPLv3 + link exception

```
usage of Dissection (C++-interface)
                                         GPLv3/CeCILL-C
   typedef dd_real quadruple; // long double for SX
   int *ptrow = new int[n + 1]; // CSR data
   int *indcol = new int[nnz];
   quadruple *coef = new quadruple[nnz];
   quadruple *x = new quadruple[n]; // solution & RHS
                                   // file to keep messages
   FILE *fp;
   DissectionSolver<quadruple> *dslv =
     new DissectionSolver<quadruple>(num_threads, true, 0, fp);
   bool isSym = false;
   bool upper_flag = false; // true for isSym = true
   bool isWhole = false; // true for isSym = true
   dslv->SymbolicFact(n, (int *)ptrow, (int *)indcol,
                      isSym, upper_flag, isWhole, ordering);
   int scaling = 1; // diagonal scaling
   double eps_pivot = 1.0e-2;
   bool kernel_detection_all = false;
   dslv->NumericFact(0, (quadruple *)coef, scaling,
                     eps_pivot, kernel_detection_all);
   bool projection = true; // solution in the image space
   bool isTrans = false; // true for A^T x = b
   bool isScaling = true; // with internal scaling
   dslv->SolveSingle(x, projection, isTrans, isScaling);
   delete dslv;
```

Hands-on tutorial : 1/7

in the frontend of CMC, login.hpc.cmc.osaka-u.ac.jp

/sc/cmc/apl/DirectSolverSeminar/ contains materials

-- MM-matrix/: matrix data from https://sparse.tamu.edu VCC/: sources, library, compiled objects ... for Intel Xeon

-- include/ header files from

MUMPS, Dissection, SCOTCH, METIS, QD, zlib lib/library files of those patch/patch files to the original distribution for compilation MUMPS/MUMPS 5.1.1 dissection/Dissection 1.0.3 metis-5.1.0/METIS 5.1.0 qd-2.3.17/QD 2.3.17 pacthed for intel compiler scotch_6.0.4/SCOTCH 6.0.4 zlib-1.2.11/Zlib 1.2.11 (used for SCOTCH) tutorial/ -- MM-pardiso.cpp, Makefile.pardiso.C++

MM-MUMPS.cpp, MM-Dissection.cpp, ..., job.pbs

Dissection/: quadruple precision solver with C++

 $\ensuremath{\texttt{SX/}}$: sources, library, compiled objects ... for NEC SX-ACE

- ...

tutorial/

Hands-on tutorial: 2/7

Makefile.Pardiso.C++

```
CXX = icpc
CCFLAGS = -03
LD = $(CXX)
LIB DIR MKL = /opt/intel/compilers and libraries 2017/linux/mkl/lib/intel64
INC DIR MKL = /opt/intel/compilers and libraries 2017/linux/mkl/include/
all: MM-Pardiso
.cpp.o:
        $(CXX) $(CCFLAGS) -I$(INC DIR MKL) -I. -c $< -o $@
MM-Pardiso: MM-Pardiso.o elapsed tme.o
        $(LD) -o MM-Pardiso -g -O3 MM-Pardiso.o elapsed time.o \
        -Xlinker -rpath=$(LIB DIR MKL) -L$(LIB DIR MKL) \
        -lmkl intel lp64 -lmkl intel thread -lmkl core -liomp5 \
       -lintlc -lsvml -lm
clean:
       rm -fr *~ *.o *.so core *.d MM-Pardiso
% mkdir VCC
% rsync -avu /sc/cmc/apl/DirectSolverSeminar/VCC/tutorial ./VCC/
% rsync -avu /sc/cmc/apl/DirectSolverSeminar/MM-matrix ./
% cd VCC/tutorial
% make -f Makefile.Pardiso.C++ all
% ldd MM-Pardiso
% qsub job.pbs
```

Hands-on tutorial: 3/7

```
% ldd MM-MUMPS
linux-vdso.so.1 => (0x00007fff8ff16000)
libdmumps.so => /sc/cmc/apl/DirectSolverSeminar/VCC/lib/libdmumps.so
libmumps common.so => /sc/cmc/apl/DirectSolverSeminar/VCC/lib/libmumps common.so
libpord.so => /sc/cmc/apl/DirectSolverSeminar/VCC/lib/libpord.so
libmpiseq.so => /sc/cmc/apl/DirectSolverSeminar/VCC/lib/libmpiseq.so
libesmumps.so => /sc/cmc/apl/DirectSolverSeminar/VCC/lib/libesmumps.so
libscotch.so => /sc/cmc/apl/DirectSolverSeminar/VCC/lib/libscotch.so
libscotcherr.so => /sc/cmc/apl/DirectSolverSeminar/VCC/lib/libscotcherr.so
libscotcherrexit.so => /sc/cmc/apl/DirectSolverSeminar/VCC/lib/libscotcherrexit.so
libmetis.so => /sc/cmc/apl/DirectSolverSeminar/VCC/lib/libmetis.so
libmkl intel lp64.so => $INTEL MKL LIB/libmkl intel lp64.so
libmkl intel thread.so => $INTEL MKL LIB/libmkl intel thread.so
libmkl core.so => $INTEL MKL LIB/libmkl core.so
libiomp5.so => $INTEL COMPILER LIB/libiomp5.so
libm.so.6 => /lib64/libm.so.6
librt.so.1 => /lib64/librt.so.1
libstdc++.so.6 => /usr/lib64/libstdc++.so.6
libgcc_s.so.1 => /lib64/libgcc_s.so.1
libc.so.6 => /lib64/libc.so.6
libdl.so.2 => /lib64/libdl.so.2
libifport.so.5 => $INTEL COMPILER LIB/libifport.so.5
libifcoremt.so.5 => $INTEL COMPILER LIB/libifcoremt.so.5
libimf.so => $INTEL COMPILER LIB/libimf.so
libsyml.so => $INTEL COMPILER LIB/libsyml.so
libintlc.so.5 => $INTEL COMPILER LIB/libintlc.so.5
libpthread.so.0 => /lib64/libpthread.so.0
libirng.so => $INTEL COMPILER LIB/libirng.so
libz.so.1 => /sc/cmc/apl/DirectSolverSeminar/VCC/lib/libz.so.1
/lib64/ld-linux-x86-64.so.2
```

INTEL_COMPILER_LIB
=/opt/intel/compilers_and_libraries_2017.2.174/linux/compiler/lib/intel64
INTEL_MKL_LIB=
/opt/intel/compilers_and_libraries_2017.2.174/linux/mkl/lib/intel64

Hands-on tutorial: 4/7

parameters for each solver

```
> Pardiso
./MM-Pardiso [matrix-market file] [num_threads] [perturbation]
iparam[1] = 2; // METIS
iparam[9] = perturbation; // 13 ↔ 10<sup>-13</sup>
```

parallelization is controlled by setting internal value of MKL

```
mkl_set_num_threads(num_threads);
```

MUMPS

parallelization is controlled by setting internal value of OpenMP

```
omp_set_dynamic(0);
omp_set_num_threads(num_threads);
```

Dissection

parallelization is controlled by initialization call

Hands-on tutorial: 6/7

3

Pardiso

error	$3.2520097 \times 10^{-14}$	
residual	$9.8931920 \times 10^{-17}$	
	CPU time (sec.)	elapsed time (sec.)
symbolic	5.6900	4.2584
numeric	87.030	21.821
fw/bw	1.7100	0.4275
MUMPS		
error	$3.1367255 \times 10^{-14}$	
rooidual	$2.1307233 \times 10^{-16}$	
residual	2.0755498 × 10	
	CPU time (sec.)	elapsed time (sec.)
symbolic	3.1900	3.1770
numeric	81.090	21.472
fw/bw	1.9400	0.4818
Dissection		
Dissection	E (D (1) E (D (D (1) E (D	
error	$7.6341570 \times 10^{-11}$	
residual	$7.1339225 \times 10^{-16}$	
	CPU time (sec.)	elapsed time (sec.)
symbolic	14.040	10.936
numeric fact	86.210	22.909
fw/bw	1.1100	0.3155

Hands-on tutorial: 7/7

results on stokes.skewsym with 4 cores : singular, dimension 2 kernel

Pardiso

error	1.6661762×10^{-2}	
residual	$1.4771827 \times 10^{-16}$	
	CPU time (sec.)	elapsed time (sec.)
symbolic	6.3400	5.3841
numeric	13.070	3.2679
fw/bw	1.0400	0.2612

MUMPS : kernel dimension detected as 2

error	$1.0595600 \times 10^{-10}$	
residual	$1.8276515 \times 10^{-14}$	
	CPU time (sec.)	elapsed time (sec.)
symbolic	4.9000	4.8968
numeric	23.120	6.2821
fw/bw	1.6600	0.4157

Dissection : kernel dimension detected as 2								
error	$7.9142093 \times 10^{-13}$							
residual	$6.2848715 \times 10^{-16}$							
	CPU time (sec.)	elapsed time (sec.)						
symbolic	12.200	9.6785						
numeric fact	18.980	5.4162						
fw/bw	0.6500	0.1840						

Web resource

Pardiso

https://software.intel.com/en-us/
mkl-developer-reference-c-intel-mkl-pardisoparallel-direct-sparse-solver-interface

https://software.intel.com/en-us/
mkl-developer-reference-fortran-intel-mkl-pardisoparallel-direct-sparse-solver-interface

MUMPS

http://mumps.enseeiht.fr

Dissection

http://www.freefem.org/ff++/ff++/download/dissection